

## 2 Problems, algorithms, and solutions

### 2.1

- (i)  $\{1\}$ ,
- (ii)  $\emptyset$ ,
- (iii)  $\mathbb{R}$ .

**2.2** Suppose that there is another solution  $x^{***}$ , say, to  $g(x) = 0$  that is different to  $x^* = -3$  and  $x^{**} = 1$ . There are three cases, depending on the relationship of  $x^{***}$  to the solutions  $x^* = -3$  and  $x^{**} = 1$ . We consider each case in turn:

(i)  $x^{***} < -3$ . Then:

$$\begin{aligned}g(x^{***}) &= (x^{***})^2 + 2x^{***} - 3, \\ &= x^{***}(x^{***} + 2) - 3, \\ &> (-3)(-1) - 3, \text{ since } x^{***} < -3, x^{***} + 2 < -1, \\ &= 0.\end{aligned}$$

(ii)  $-3 < x^{***} < 1$ . Then:

$$\begin{aligned}g(x^{***}) &= (x^{***})^2 + 2x^{***} - 3, \\ &= (x^{***} + 1)^2 - 4, \\ &< (2)^2 - 4, \text{ since } |x^{***} + 1| < 2, \\ &= 0.\end{aligned}$$

(iii)  $x^{***} > 1$ . Then:

$$\begin{aligned}g(x^{***}) &= (x^{***})^2 + 2x^{***} - 3, \\ &= x^{***}(x^{***} + 2) - 3, \\ &> (1)(3) - 3, \text{ since } x^{***} > 1, x^{***} + 2 > 3, \\ &= 0.\end{aligned}$$

In each case,  $g(x^{***}) \neq 0$ , so no such solution  $x^{***}$  exists that is different to  $x^* = -3$  and  $x^{**} = 1$ .

### 2.3

- (i) 1,

(ii)  $\{2\}$ .

**2.4** Suppose that  $\underline{f} \leq 1$ . Then:

$$\begin{aligned}\underline{f} &\leq 1, \\ &\leq (x-2)^2 + 1, \forall x \in \mathbb{R}, \text{ since } (x-2)^2 \geq 0.\end{aligned}$$

So,  $\underline{f}$  is a lower bound for the problem  $\min_{x \in \mathbb{S}} f(x)$  according to Definition 2.2.

**2.5** Suppose that  $\underline{f} \leq f^*$ . Then:

$$\begin{aligned}\underline{f} &\leq f^*, \\ &\leq f(x), \forall x \in \mathbb{S},\end{aligned}$$

by definition of minimum. That is,  $\underline{f}$  is a lower bound for  $\min_{x \in \mathbb{S}} f(x)$  according to Definition 2.2.

**2.6**

Part	$x^*$	$x^{**}$	$x^{***}$
(i) $h_1(x) \leq 0$ active?	Yes	Yes	No
(ii) $h_2(x) \leq 0$ active?	Yes	No	No
(iii) Active set?	$\{1, 2\}$	$\{1\}$	$\emptyset$
(iv) Strictly feasible for $h_1(x) \leq 0$ ?	No	No	Yes
(v) Strictly feasible for $h_2(x) \leq 0$ ?	No	Yes	Yes
(vi) Strictly feasible for $h(x) \leq \mathbf{0}$ ?	No	No	Yes
(vii) On boundary of $\{x \in \mathbb{R}^2   h(x) \leq \mathbf{0}\}$ ?	Yes	Yes	No

**2.7**

(i) The contour set is defined by:

$$\begin{aligned}\mathbb{C}_f(\tilde{f}) &= \{x \in \mathbb{S} | f(x) = \tilde{f}\}, \\ &= \{x \in \mathbb{S} | (x_1)^2 + (x_2 + 1)^2 - 4 = \tilde{f}\},\end{aligned}$$

which is the set of points  $\begin{bmatrix} x_1 \\ x_2 \end{bmatrix}$  on the circle of radius  $\sqrt{\tilde{f} + 4}$  and center

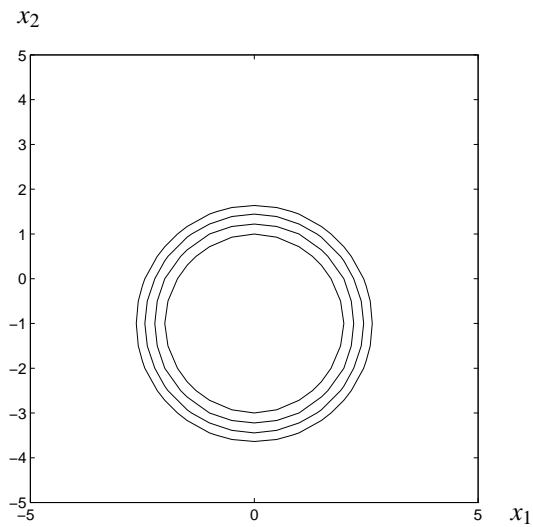


Fig. 1. Contour sets of function in Exercise 2.7, Part (i).

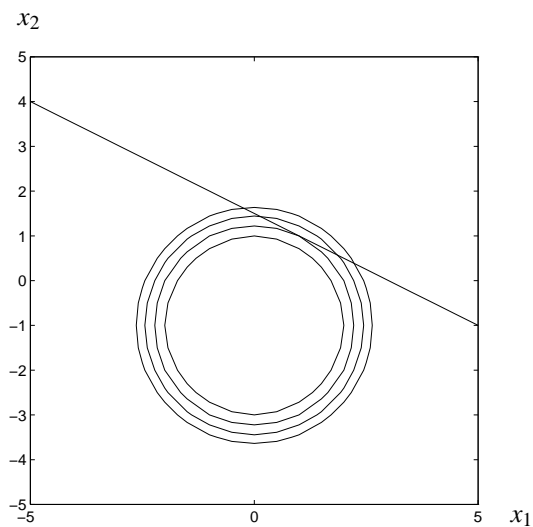


Fig. 2. Points satisfying  $g(x) = \mathbf{0}$  and contour sets for Exercise 2.7, Part (ii).

$\begin{bmatrix} 0 \\ -1 \end{bmatrix}$ . For  $\tilde{f} = 0, 1, 2, 3$ , the contour sets are the circles of radius  $2, \sqrt{5}, \sqrt{6}$ , and  $\sqrt{7}$ , respectively, all with center  $\begin{bmatrix} 0 \\ -1 \end{bmatrix}$ . The contour sets are shown in Figure 1.

- (ii) The set of points satisfying  $g(x) = \mathbf{0}$  along with the contour sets of the function  $f$  are shown in Figure 2.

(iii) From Figure 2,

$$\begin{aligned}\min_{x \in \mathbb{R}^2} \{f(x) | x_1 + 2x_2 - 3 = 0\} &= 1, \\ \operatorname{argmin}_{x \in \mathbb{R}^2} \{f(x) | x_1 + 2x_2 - 3 = 0\} &= \left\{ \begin{bmatrix} 1 \\ 1 \end{bmatrix} \right\}.\end{aligned}$$

## 2.8

(i) The variables are defined as follows:

- $x_1$  is the bandwidth for customer 1;
- $x_2$  is the bandwidth for customer 2; and
- $x_3$  is the bandwidth for customer 3.

Note that customer 3 only cares about getting data from point X to point Z, so there is only one variable associated with customer 3. We collect the three entries together into a vector  $x \in \mathbb{R}^3$ .

(ii) The objective is  $f : \mathbb{R}^3 \rightarrow \mathbb{R}$  defined by:

$$\forall x \in \mathbb{R}_+^3, f(x) = \sum_{k=1}^3 f_k(x_k).$$

(iii) With these definitions of variables, there are no equality constraints.

(iv) Both customer 1 and customer 3 traffic requires bandwidth on link a and the maximum bandwidth on this link is  $c_a$ . Both customer 2 and customer 3 traffic requires bandwidth on link b and the maximum bandwidth on this link is  $c_b$ . Since bandwidth must be positive and also since the objective function is only defined for non-negative values of bandwidth, we must also incorporate non-negativity constraints. Therefore, we can express the inequality constraints in the form  $Cx \leq d$  with:

$$\begin{aligned}C &= \begin{bmatrix} 1 & 0 & 1 \\ 0 & 1 & 1 \\ -1 & 0 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \end{bmatrix}, \\ d &= \begin{bmatrix} c_a \\ c_b \\ 0 \\ 0 \\ 0 \end{bmatrix}.\end{aligned}$$

If the functions  $f_k$  were known to be strictly monotonically increasing then an alternative formulation would be to represent the capacity constraints as equalities since all capacity would be used at the optimum.

## 2.9

(i) Since  $aB - Ab \neq 0$  we have unique solution:

$$\begin{bmatrix} X^* \\ Y^* \end{bmatrix} = \frac{1}{aB - Ab} \begin{bmatrix} B & -b \\ -A & a \end{bmatrix} \begin{bmatrix} c \\ C \end{bmatrix}.$$

(ii) The conditions are that:

$$\begin{aligned} aB - Ab &\neq 0, \\ \frac{1}{aB - Ab}(Bc - bC) &\geq 0, \\ \frac{1}{aB - Ab}(-Ac + aC) &\geq 0. \end{aligned}$$

(iii) If the conditions in the previous part are satisfied, then the solutions are the four possible values:

$$\begin{bmatrix} x^* \\ y^* \end{bmatrix} = \begin{bmatrix} \sqrt{\frac{1}{aB - Ab}(Bc - bC)} \\ \sqrt{\frac{1}{aB - Ab}(-Ac + aC)} \end{bmatrix}, \begin{bmatrix} \sqrt{\frac{1}{aB - Ab}(Bc - bC)} \\ -\sqrt{\frac{1}{aB - Ab}(-Ac + aC)} \end{bmatrix}, \\ \begin{bmatrix} -\sqrt{\frac{1}{aB - Ab}(Bc - bC)} \\ \sqrt{\frac{1}{aB - Ab}(-Ac + aC)} \end{bmatrix}, \begin{bmatrix} -\sqrt{\frac{1}{aB - Ab}(Bc - bC)} \\ -\sqrt{\frac{1}{aB - Ab}(-Ac + aC)} \end{bmatrix}.$$

We could write this more compactly as:

$$\begin{bmatrix} x^* \\ y^* \end{bmatrix} = \begin{bmatrix} \pm \sqrt{\frac{1}{aB - Ab}(Bc - bC)} \\ \pm \sqrt{\frac{1}{aB - Ab}(-Ac + aC)} \end{bmatrix}.$$

However, if this notation is used then, to avoid confusion, it is important to explicitly note that all four combinations of plus and minus are valid solutions.

**2.10** Suppose that an algorithm did exist for finding the minimum and all minimizers of unconstrained minimization problems  $\min_{x \in \mathbb{R}^n} f(x)$ , where  $f: \mathbb{R}^n \rightarrow \mathbb{R}$  is an arbitrary partially differentiable function. Let  $g: \mathbb{R} \rightarrow \mathbb{R}$  be an arbitrary polynomial in a single variable and define  $f: \mathbb{R} \rightarrow \mathbb{R}$  by:

$$\forall x \in \mathbb{R}, f(x) = (g(x))^2.$$

The function  $f$  is partially differentiable since  $g$  is a polynomial, which is partially differentiable. Note that if  $\min_{x \in \mathbb{R}^n} f(x) \neq 0$  then there is no solution to  $g(x) = 0$ . On the other hand, if  $\min_{x \in \mathbb{R}^n} f(x) = 0$  and  $x^* \in \operatorname{argmin}_{x \in \mathbb{R}^n} f(x)$  then  $g(x^*) = 0$  and the minimizers of  $\min_{x \in \mathbb{R}^n} f(x) = 0$  are the solutions of  $g(x) = 0$ . That is, the existence of a direct algorithm to minimize  $f$  and find all of its minimizers would also yield a direct algorithm to solve  $g(x) = 0$ . But no such direct algorithm to solve  $g(x) = 0$  exists, so no such algorithm to minimize  $f$  and find all of its minimizers can exist.

**2.11** Let  $\|\bullet\|$  be the absolute value norm and let  $\varepsilon > 0$  be given. We claim that  $N = \lceil 1/\varepsilon \rceil$  will suffice in Definition 2.9, where  $\lceil \bullet \rceil$  is the smallest integer that is greater than or equal to its argument. For, let  $v \geq N$ . Then

$$\begin{aligned} \|x^{(v)} - 0\| &= \left| \frac{1}{v+1} \right|, \\ &= \frac{1}{v+1}, \\ &\leq \frac{1}{N+1}, \\ &\leq \frac{1}{N}, \\ &\leq \varepsilon, \end{aligned}$$

since  $N \geq 1/\varepsilon$ .

## 2.12

(i)

$$\begin{aligned} \forall v \in \mathbb{Z}_+, \|x^{(v+1)} - x^*\| &\leq C \|x^{(v)} - x^*\|, \\ &\leq (C)^2 \|x^{(v-1)} - x^*\|, \\ &\vdots \\ &\leq (C)^{v+1} \|x^{(0)} - x^*\|, \\ &\leq (C)^{v+1} \bar{\rho}. \end{aligned}$$

So,  $\|x^{(N)} - x^*\| \leq (C)^N \bar{\rho} \leq \varepsilon \bar{\rho}$  if  $(C)^N \leq \varepsilon$ , which is true if:

$$N \ln(C) \leq \ln(\varepsilon),$$

or  $N \geq \ln(\varepsilon)/\ln(C)$ , noting that  $0 < C < 1$ . That is,  $N = \lceil \ln(\varepsilon)/\ln(C) \rceil$ .

(ii) We first note that:

$$\begin{aligned} ((\alpha)^2)^2 &= (\alpha)^4, \\ &= (\alpha)^{(2)^2}, \\ (((\alpha)^2)^2)^2 &= ((\alpha)^{(2)^2})^2, \\ &= (\alpha)^{(2)^3}, \\ (\dots(\alpha)^2 \dots)^2 &= (\alpha)^{(2)^v}, \end{aligned}$$

where there are  $v$  exponentiations in total in the left-hand side of the last line. Therefore:

$$\begin{aligned} \forall v \in \mathbb{Z}_+, \left\| x^{(v+1)} - x^* \right\| &\leq C \left\| x^{(v)} - x^* \right\|^2, \\ &\leq C(C)^2 \left\| x^{(v-1)} - x^* \right\|^4, \\ &= C(C)^2 \left\| x^{(v-1)} - x^* \right\|^{(2)^2}, \\ &\leq C(C)^2 (C)^{(2)^2} \left\| x^{(v-2)} - x^* \right\|^8, \\ &= C(C)^2 (C)^{(2)^2} \left\| x^{(v-2)} - x^* \right\|^{(2)^3}, \\ &\vdots \\ &\leq C(C)^2 (C)^{(2)^2} \dots (C)^{(2)^v} \left\| x^{(0)} - x^* \right\|^{(2)^{(v+1)}}, \\ &= (C)^{(2)^{v+1}-1} \left\| x^{(0)} - x^* \right\|^{(2)^{(v+1)}}, \\ &\leq (C)^{(2)^{v+1}-1} (\bar{\rho})^{(2)^{(v+1)}}. \end{aligned}$$

So,  $\left\| x^{(N)} - x^* \right\| \leq (C)^{(2)^N-1} (\bar{\rho})^{(2)^N} \leq \varepsilon \bar{\rho}$  if  $(C)^{(2)^N-1} (\bar{\rho})^{(2)^N-1} \leq \varepsilon$ , which is true if:

$$((2)^N - 1) \ln(C) + ((2)^N - 1) \ln(\bar{\rho}) \leq \ln(\varepsilon),$$

or  $((2)^N - 1)(\ln(C) + \ln(\bar{\rho})) \leq \ln(\varepsilon)$ . Since  $\varepsilon < 1$  and  $\ln(\varepsilon) < 0$  this inequality can only be true if  $\ln(C) + \ln(\bar{\rho}) < 0$ , which means that we must require  $C\bar{\rho} < 1$ . If this is true, then the condition becomes:

$$(2)^N - 1 \geq \frac{\ln(\varepsilon)}{\ln(C) + \ln(\bar{\rho})},$$

or:

$$N \geq \ln \left( \frac{\ln(\varepsilon)}{\ln(C) + \ln(\bar{\rho})} + 1 \right) / \ln(2).$$

That is:

$$N = \left\lceil \ln \left( \frac{\ln(\varepsilon)}{\ln(C) + \ln(\bar{\rho})} + 1 \right) / \ln(2) \right\rceil.$$

(iii) We first note that:

$$\begin{aligned} ((\alpha)^R)^R &= (\alpha)^{((R)^2)}, \\ (((\alpha)^R)^R)^R &= (((\alpha)^{((R)^2)})^R), \\ &= (\alpha)^{((R)^3)}, \\ (\dots (\alpha)^R \dots)^R &= (\alpha)^{((R)^v)}, \end{aligned}$$

where there are  $v$  exponentiations in total in the left-hand side of the last line. Therefore:

$$\begin{aligned} \forall v \in \mathbb{Z}_+, \left\| x^{(v+1)} - x^* \right\| &\leq C \left\| x^{(v)} - x^* \right\|^R, \\ &\leq C(C)^R \left\| x^{(v-1)} - x^* \right\|^{((R)^2)}, \\ &\leq C(C)^R (C)^{((R)^2)} \left\| x^{(v-2)} - x^* \right\|^{((R)^3)}, \\ &\vdots \\ &\leq C(C)^R (C)^{((R)^2)} \dots (C)^{(R^v)} \left\| x^{(0)} - x^* \right\|^{((R)^{(v+1)})}, \\ &= (C)^{\frac{((R)^{v+1}-1)}{(R-1)}} \left\| x^{(0)} - x^* \right\|^{((R)^{(v+1)})}, \\ &\leq (C)^{\frac{((R)^{v+1}-1)}{(R-1)}} (\bar{\rho})^{((R)^{(v+1)})}. \end{aligned}$$

So,  $\left\| x^{(N)} - x^* \right\| \leq (C)^{\frac{((R)^N-1)}{(R-1)}} (\bar{\rho})^{((R)^N)} \leq \varepsilon \bar{\rho}$  if  $(C)^{\frac{((R)^N-1)}{(R-1)}} (\bar{\rho})^{((R)^N-1)} \leq \varepsilon$ , which is true if:

$$\frac{((R)^N-1)}{(R-1)} \ln(C) + ((R)^N-1) \ln(\bar{\rho}) \leq \ln(\varepsilon),$$

or  $((R)^N-1) \left( \frac{\ln(C)}{R-1} + \ln(\bar{\rho}) \right) \leq \ln(\varepsilon)$ . Since  $\varepsilon < 1$  and  $\ln(\varepsilon) < 0$  this inequality can only be true if  $\ln(C)/(R-1) + \ln(\bar{\rho}) < 0$ , which means that we must require  $(C)^{1/(R-1)} \bar{\rho} < 1$ . If this is true, then the condition becomes:

$$(R)^N - 1 \geq \frac{\ln(\varepsilon)}{\frac{\ln(C)}{R-1} + \ln(\bar{\rho})},$$

or:

$$N \geq \ln \left( \frac{\ln(\varepsilon)}{\frac{\ln(C)}{R-1} + \ln(\bar{\rho})} + 1 \right) / \ln(R).$$

That is:

$$N = \left\lceil \ln \left( \frac{\ln(\varepsilon)}{\frac{\ln(C)}{R-1} + \ln(\bar{\rho})} + 1 \right) / \ln(R) \right\rceil.$$

**2.13** We consider each sequence in turn and evaluate the corresponding limit.

(i)  $\forall v \in \mathbb{Z}_+, x^{(v)} = 1/(v+1)$ .

(a)  $R = 0$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|1/(v+2)\|}{\|1/(v+1)\|^0}, \\ &= \lim_{v \rightarrow \infty} \|1/(v+2)\|, \\ &= 0. \end{aligned}$$

(b)  $R = \frac{1}{2}$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|1/(v+2)\|}{\|1/(v+1)\|^{1/2}}, \\ &= \lim_{v \rightarrow \infty} \frac{\sqrt{v+1}}{v+2}, \\ &= \lim_{v \rightarrow \infty} \frac{\sqrt{v+1}}{\sqrt{v+2}} \frac{1}{\sqrt{v+2}}, \\ &\leq \lim_{v \rightarrow \infty} \frac{1}{\sqrt{v+2}}, \\ &= 0. \end{aligned}$$

(c)  $R = 1$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|1/(v+2)\|}{\|1/(v+1)\|}, \\ &= \lim_{v \rightarrow \infty} \frac{v+1}{v+2}, \\ &= 1. \end{aligned}$$

(d)  $R = 2$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|1/(v+2)\|}{\|1/(v+1)\|^2}, \\ &= \lim_{v \rightarrow \infty} \frac{(v+1)^2}{(v+2)}, \\ &= \infty. \end{aligned}$$

(ii)  $\forall v \in \mathbb{Z}_+, x^{(v)} = (2)^{-v}$ .

(a)  $R = 0$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-v-1}\|}{\|(2)^{-v}\|^0}, \\ &= \lim_{v \rightarrow \infty} \|(2)^{-v-1}\|, \\ &= 0. \end{aligned}$$

(b)  $R = \frac{1}{2}$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-v-1}\|}{\|(2)^{-v}\|^{1/2}}, \\ &= \lim_{v \rightarrow \infty} (2)^{-(v/2)-1}, \\ &= 0. \end{aligned}$$

(c)  $R = 1$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-v-1}\|}{\|(2)^{-v}\|}, \\ &= \lim_{v \rightarrow \infty} (2)^{-1}, \\ &= 1/2. \end{aligned}$$

(d)  $R = 2$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-v-1}\|}{\|(2)^{-v}\|^2}, \\ &= \lim_{v \rightarrow \infty} (2)^{v-1}, \\ &= \infty. \end{aligned}$$

(iii)  $\forall v \in \mathbb{Z}_+, x^{(v)} = (2)^{-((2)^v)}$ .

(a)  $R = 0$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-((2)^{v+1})}\|}{\|(2)^{-((2)^v)}\|^0}, \\ &= \lim_{v \rightarrow \infty} \|(2)^{-((2)^{v+1})}\|, \\ &= 0. \end{aligned}$$

(b)  $R = \frac{1}{2}$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-((2)^{v+1})}\|}{\|(2)^{-((2)^v)}\|^{1/2}}, \\ &= \lim_{v \rightarrow \infty} (2)^{-(3(2)^{v-1})}, \\ &= 0. \end{aligned}$$

(c)  $R = 1$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-((2)^{v+1})}\|}{\|(2)^{-((2)^v)}\|}, \\ &= \lim_{v \rightarrow \infty} (2)^{-((2)^v)}, \\ &= 0. \end{aligned}$$

(d)  $R = 2$ :

$$\begin{aligned} \lim_{v \rightarrow \infty} \frac{\|x^{(v+1)}\|}{\|x^{(v)} - x^*\|^R} &= \lim_{v \rightarrow \infty} \frac{\|(2)^{-((2)^{v+1})}\|}{\|(2)^{-((2)^v)}\|^2}, \\ &= \lim_{v \rightarrow \infty} 1, \\ &= 1. \end{aligned}$$

The limits are shown in the following table along with the rate of convergence.

$x^{(v)}$	$R =$				Rate
	0	1/2	1	2	
$1/(v+1)$	0	0	1	$\infty$	1 (but neither linear nor quadratic convergence)
$(2)^{-v}$	0	0	$\frac{1}{2}$	$\infty$	1 (linear convergence)
$(2)^{-((2)^v)}$	0	0	0	1	2 (quadratic convergence)

**2.14** Let  $x \in \mathbb{S}$  and assume that  $f(x) \geq f(x^*)$ . Define  $\phi : [0, 1] \rightarrow \mathbb{R}$  by:

$$\forall t \in [0, 1], \phi(t) = f(x^* + t(x - x^*)).$$

We have that:

$$\begin{aligned}\forall t \in [0, 1], \frac{d\phi}{dt}(t) &= \nabla f(x^* + t(x - x^*))^\dagger (x - x^*), \\ \frac{d\phi}{dt}(0) &= \nabla f(x^*)^\dagger (x - x^*), \\ \frac{d^2\phi}{dt^2}(t) &= (x - x^*)^\dagger \nabla^2 f(x^* + t(x - x^*)) (x - x^*).\end{aligned}$$

To bound the rate of convergence and rate constant, we prove two inequalities. To prove the first inequality, note that:

$$\begin{aligned}f(x) - f(x^*) &= \phi(1) - \phi(0), \\ &= \int_{t=0}^1 \frac{d\phi}{dt}(t) dt, \\ &\quad \text{by the fundamental theorem of integral calculus applied to } \phi, \\ &\quad \text{(see Theorem A.2 in Section A.4.4.1 of Appendix A),} \\ &= \int_{t=0}^1 \nabla f(x^* + t(x - x^*))^\dagger (x - x^*) dt, \\ &\leq \int_{t=0}^1 \|\nabla f(x^* + t(x - x^*))\| \|x - x^*\| dt, \\ &\leq \int_{t=0}^1 \bar{\kappa} \|x - x^*\| dt, \text{ by assumption on } \nabla f, \\ &= \bar{\kappa} \|x - x^*\|.\end{aligned}$$

Since  $f(x^{(v+1)}) \geq f(x^*)$ , we have that:

$$\begin{aligned}\|f(x^{(v+1)}) - f(x^*)\| &= f(x^{(v+1)}) - f(x^*), \\ &\leq \bar{\kappa} \|x^{(v+1)} - x^*\|.\end{aligned}$$